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# Characterising knotting properties of polymers in nanochannels†

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Using a lattice model of polymers in a tube, we define one way to characterise different configurations of a given knot as either “local” or “non-local”, based on a standard approach for measuring the “size” of a knot within a knotted polymer chain. The method involves associating knot-types to subarcs of the chain, and then identifying a knotted subarc with minimal arclength; this arclength is then the knot-size. If the resulting knot-size is small relative to the whole length of the chain, then the knot is considered to be localised or “local”; otherwise, it is “non-local”. Using this definition, we establish that all but exponentially few sufficiently long self-avoiding polygons (closed chains) in a tubular sublattice of the simple cubic lattice are “non-locally” knotted. This is shown to also hold for the case when the same polygons are subject to an external tensile force, as well as in the extreme case when they are as compact as possible (no empty lattice sites). We also provide numerical evidence for small tube sizes that at equilibrium non-local knotting is more likely than local knotting, regardless of the strength of the stretching or compressing force. The relevance of these results to other models and recent experiments involving DNA knots is also discussed.

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## 1 Introduction

Motivated in part by experimental studies of DNA packing in viral capsids<sup>1,2</sup> and DNA translocation through nanopores,<sup>3–6</sup> there has been much recent interest in understanding and characterising the entanglement complexity of confined polymers, and determining any dependencies on the extent or the mechanism of confinement. For knots in polymers, one measure of interest has been the average “size” of the knotted part of the polymer. With such a measurement, one can then characterise the knotting as “local” when the size of the knotted part is small compared to the whole length of the polymer, or otherwise as “non-local”. Based on polymer scaling theory and supporting numerical evidence, it is generally accepted that local knotting is dominant for unconfined polymers, as polymer length grows, while computer simulation studies of knotting in collapsed or spherically confined polymers suggest that non-local knotting dominates.<sup>7</sup> Recent simulation and experimental studies<sup>3,5,6,8</sup> of an intermediate regime of confinement, where the polymer is confined to a channel, tube or pore, such that polymer growth or motion is tightly restricted in two spatial dimensions but unrestricted in the third, has indicated that local

knotting dominates, unlike when all three spatial dimensions are restricted. Some differences between the entanglement complexity of open *versus* closed chains were also observed.

In this paper we explore this latter type of confinement further using lattice models of both closed (self-avoiding polygons) and open (self-avoiding walk) chains in tubular subsets of the simple cubic lattice. One major advantage of the lattice model approach is that we are able to prove results related to knot localisation for arbitrary tube dimensions and, for the case of small tube sizes, we are able to perform exact calculations related to this. We study an equilibrium model of polymers in a tube subject to a tensile force  $f$  and prove results about the limiting free energy and the likelihood of occurrence of different knotted patterns as a function of polymer length.

Our conclusions are presented as a series of eight Results. Specifically, we provide both theoretical and numerical evidence that for a knotted ring polymer at equilibrium in a nanochannel or nanopore, a knot configuration such as that shown in Fig. 1(a) is more likely than that shown in Fig. 1(b), regardless of the strength of the force  $f$  and whether or not it is a compressing ( $f < 0$ ) or a stretching ( $f > 0$ ) force (see Results 1–2). Furthermore, knot configurations such as that in Fig. 1(a) are on average tighter than those in Fig. 1(b) and hence may be expected to translocate through a nanopore faster (Result 3). We also establish for the lattice model that, as observed experimentally, the situation is quite different for open chains, with knot configurations like that of Fig. 1(e) being rare (in fact exponentially rare) compared to that shown in Fig. 1(f), because of the significant

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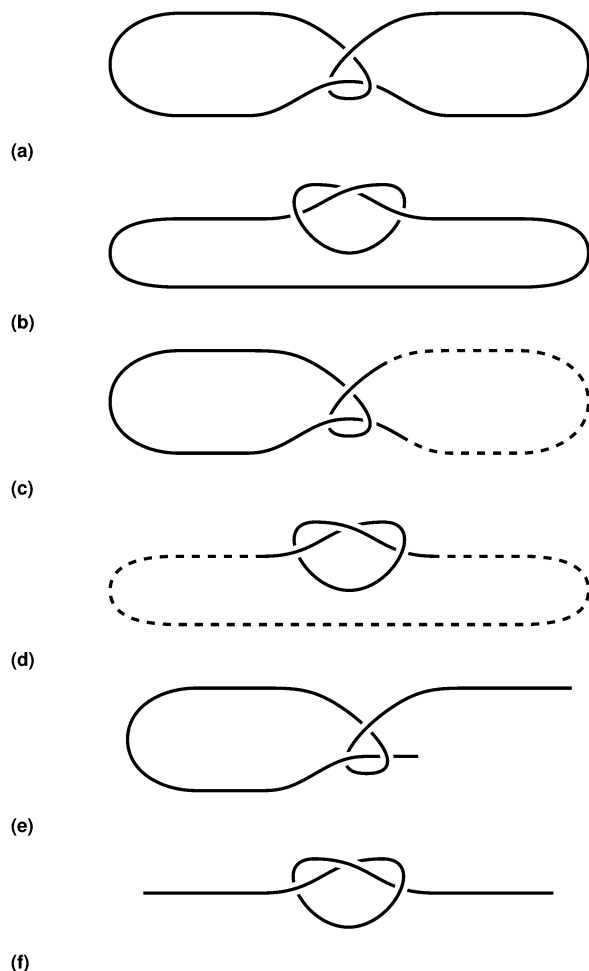


Fig. 1 Illustrations of (a) a non-local trefoil and (b) a local trefoil, with knotted subarcs drawn with solid lines in (c) and (d); open chains with (e) a non-local trefoil component and (f) a local trefoil component.

entropic disadvantage associated with the formation of a long bend (Result 8).

To obtain and explain these results, it is necessary to first characterise the differences between the different knotted configurations in Fig. 1. We do this here by using two different measures of knot size. For example, when one measures the size of the knotted part according to the size of the region in which the crossings are concentrated, then the configurations shown in both Fig. 1(a) and (b) correspond to examples of “tight” knotting and this is consistent, for example, with a measure of how long it would take the knotted part to translocate through a nanopore. However, using another standard measure for determining knot-size, namely using the length of a smallest knotted subarc, leads to characterising Fig. 1(a) as “non-local” knotting (a knotted subarc is drawn with a solid line in Fig. 1(c)); while that in Fig. 1(b) is “local” (a knotted subarc is drawn with a solid line in Fig. 1(d)). Using this latter measure, our results indicate that non-local knotting is more likely than local knotting in a tube. This suggests that the likelihood of non-local knotting in other restricted spaces, like spheres, will be dependent on the definition chosen for knot-size.

We note that an existing model of DNA in a nanochannel (Micheletti and Orlandini<sup>3</sup>) considers DNA in a salt solution; a corresponding lattice model would include long-range interactions to take into account screened Coulomb interactions such as in Tesi *et al.*<sup>9</sup> Including such long range interactions increases the complexity of the model and makes it less tractable to both theoretical and numerical analysis. Hence, as a first step, we explore here a model where we can study a range of more tractable scenarios which include the good solvent regime as well as the fully compressed regime. Because of this, we do not make a direct connection to DNA experiments. However, we do argue that the tube sizes for which we have numerical results correspond to tube dimensions within the 10–40 nm range and where the polymers are well below the de Gennes scaling region identified in Micheletti and Orlandini.<sup>3</sup>

The remainder of this paper is organised as follows. In Section 2 we present definitions of the models under consideration. In Section 3, definitions for two measures of knot size and our classification scheme for knotted patterns are given. After that, we present exact and numerical results about the models. Then we present our theoretical results (outlines of proofs are given in the ESI†) and the methods used. In Section 6 we briefly discuss some other models of confined DNA and contrast and compare with the models studied here.

## 2 The models

We will use a general model for polygons in lattice tubes subject to an external force which has been studied previously; the notation and definitions used here (unless stated otherwise) are as in ref. 10.

For non-negative integers  $L, M$ , let  $\mathbb{T}_{L,M} \equiv \mathbb{T} \subset \mathbb{Z}^3$  be the semi-infinite  $L \times M$  tube on the cubic lattice defined by

$$\mathbb{T} = \{(x,y,z) \in \mathbb{Z}^3 : x \geq 0, 0 \leq y \leq L, 0 \leq z \leq M\}.$$

Define  $\mathcal{P}_{\mathbb{T}}$  to be the set of self-avoiding polygons in  $\mathbb{T}$  which occupy at least one vertex in the plane  $x = 0$ , and let  $\mathcal{P}_{\mathbb{T},n}$  be the subset of  $\mathcal{P}_{\mathbb{T}}$  comprising polygons with  $n$  edges ( $n$  even). Then let  $p_{\mathbb{T},n} = |\mathcal{P}_{\mathbb{T},n}|$ .

We define the span  $s(\pi)$  of a polygon  $\pi \in \mathcal{P}_{\mathbb{T}}$  to be the maximal  $x$ -coordinate reached by any of its vertices and we use  $|\pi|$  to denote the number of edges in  $\pi$ . See Fig. 2(a) for a polygon  $\pi$  that fits in a  $2 \times 1$  tube with  $s(\pi) = 6$  and  $|\pi| = 36$ . To model a force acting parallel to the  $x$ -axis, we associate a fugacity (Boltzmann weight)  $e^{fs(\pi)}$  with each polygon  $\pi$ . Let  $p_{\mathbb{T},n}(s)$  be the number of polygons in  $\mathcal{P}_{\mathbb{T},n}$  with span  $s$ . Then the “fixed-edge” model partition function is given by

$$Z_{\mathbb{T},n}(f) = \sum_{|\pi|=n} e^{fs(\pi)} = \sum_s p_{\mathbb{T},n}(s) e^{fs}.$$

Thus  $f \ll 0$  corresponds to the “compressed” regime while  $f \gg 0$  corresponds to the “stretched” regime. For this model, the probability of a polygon  $\pi \in \mathcal{P}_{\mathbb{T},n}$  is given by

$$\mathbb{P}_n^{(\text{ed},f)}(\pi) = \frac{e^{fs(\pi)}}{Z_{\mathbb{T},n}(f)}.$$

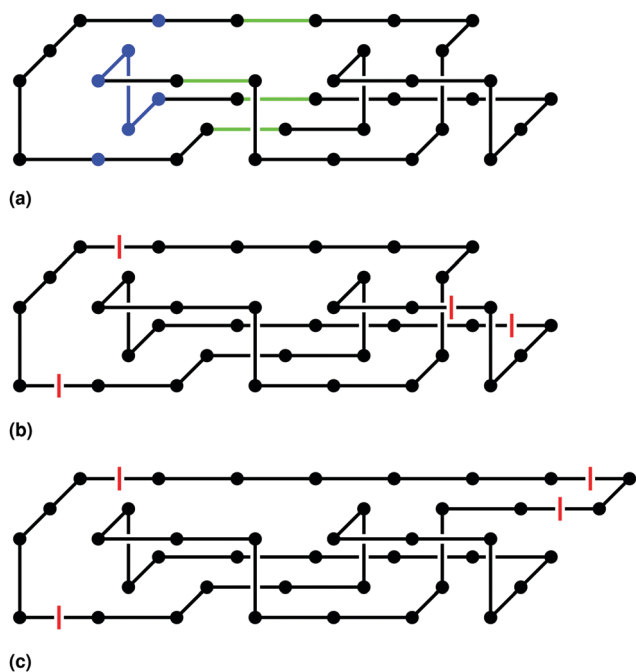


Fig. 2 (a) A 36-edge polygon  $\pi$  that fits inside  $\mathbb{T}_{L,M}$  with  $L \geq 2$  and  $M \geq 1$ ; the tube extends without bound to the right and the span  $s(\pi) = 6$ . Blue vertices and edges denote a hinge of  $\pi$ , and green edges denote a section of  $\pi$ . (b) The locations of the two pairs of vertical red lines indicate the locations of the two 2-sections in this polygon; in this example, the polygon can be decomposed into a start unknot pattern, a proper trefoil knot pattern, and an end unknot pattern. The proper knot pattern is classified as non-local in this case. (c) A local proper knot pattern in the same tube with span 7.

The (limiting) free energy per edge of polygons in  $\mathbb{T}$  is defined as

$$\mathcal{F}_{\mathbb{T}}(f) = \lim_{n \rightarrow \infty} \frac{1}{n} \log Z_{\mathbb{T},n}(f).$$

This is known<sup>11</sup> to exist for all  $f$ .

For  $f = 0$ , it has been proved that<sup>12,13</sup>

$$\begin{aligned} \mathcal{F}_{\mathbb{T}}(0) &= \lim_{n \rightarrow \infty} n^{-1} \log p_{\mathbb{T},n} \\ &< \lim_{n \rightarrow \infty} n^{-1} \log c_{\mathbb{T},n} \\ &< \lim_{n \rightarrow \infty} n^{-1} \log p_n \\ &= \lim_{n \rightarrow \infty} n^{-1} \log c_n \equiv \kappa \equiv \log \mu, \end{aligned} \quad (1)$$

where  $c_n$  is the number of  $n$ -step self-avoiding walks (SAWs) in  $\mathbb{Z}^3$  starting at the origin and  $\kappa$  is their connective constant, and  $c_{\mathbb{T},n}$  is the number of these confined to  $\mathbb{T}$ .

A subset of self-avoiding polygons in  $\mathbb{T}$  are Hamiltonian polygons: those which occupy every vertex in a  $s \times L \times M$  subtube of  $\mathbb{T}$ . These serve as an idealised model of tightly packed ring polymers, in addition to being a useful lower bound for general polygons in the  $f < 0$  compressed regime. We define the number of Hamiltonian polygons,  $p_{\mathbb{T},n}^H$ , to be the number of  $n$ -edge polygons in  $\mathcal{P}_{\mathbb{T},n}$  which have span  $s$  and occupy every vertex in

an  $s \times L \times M$  subtube of  $\mathbb{T}$ . We define  $W = (L + 1)(M + 1)$  (the number of vertices in an integer plane  $x = i \geq 0$  of the tube) and will assume without loss of generality that  $L \geq M$ ; note that  $p_{\mathbb{T},n}^H = 0$  if  $n$  is not a multiple of  $W$ . The following limit has been proved to exist by Beaton *et al.*<sup>10</sup> (see also Eng<sup>14</sup>):

$$\kappa_{\mathbb{T}}^H \equiv \lim_{s \rightarrow \infty} \frac{1}{(s+1)W} \log p_{\mathbb{T},(s+1)W}^H.$$

Furthermore, using this,  $\mathcal{F}_{\mathbb{T}}(f)$ , the free energy per edge, is bounded as follows:

$$\max\{f/2, (fW) + \kappa_{\mathbb{T}}^H\} \leq \mathcal{F}_{\mathbb{T}}(f) \leq \max\{fW, f/2\} + \mathcal{F}_{\mathbb{T}}(0), \quad (2)$$

with  $\mathcal{F}_{\mathbb{T}}(f)$  asymptotic to the lower bound for  $f \rightarrow \infty$  for any  $\mathbb{T}$ , and for  $f \rightarrow -\infty$  for small tube sizes (this is conjectured to be true for any  $\mathbb{T}$ ), see Beaton *et al.*<sup>10</sup>

Here, we will also be interested in the dual model, called the “fixed-span” model, with partition function given by

$$Q_{\mathbb{T},s}(g) = \sum_n p_{\mathbb{T},n}(s) e^{gn}.$$

For this partition function, when  $g \gg 0$  densely packed (in terms of number of edges per span) polygons dominate the partition function, while when  $g \ll 0$  polygons with very few edges per span dominate. For this model, the probability associated with a span  $s$  polygon  $\pi$  is given by

$$\mathbb{P}_s^{(\text{sp},g)}(\pi) = \frac{e^{g|\pi|}}{Q_{\mathbb{T},s}(g)}$$

and the associated (limiting) free energy per span exists<sup>10</sup> (see also ref. 15):

$$\mathcal{G}_{\mathbb{T}}(g) = \lim_{s \rightarrow \infty} \frac{1}{s} \log Q_{\mathbb{T},s}(g).$$

Both models correspond to special cases of the grand canonical partition function

$$G_{\mathbb{T}}(f, g) = \sum_s \sum_n p_{\mathbb{T},n}(s) e^{gn+fs}$$

and can be studied using transfer matrix methods.<sup>15</sup>

Hamiltonian polygons can also be studied using transfer matrices<sup>10,14</sup> and we will also be interested in the fixed-span model where polygons are restricted to being Hamiltonian and the probability associated with a span  $s$  Hamiltonian polygon  $\pi$  is given by

$$\mathbb{P}_s^{(H)}(\pi) = \frac{1}{p_{\mathbb{T},(s+1)W}^H}.$$

### 3 Characterising local knotting and classifying knotted patterns

Given a polygon of prime knot-type  $K$ , one standard approach for measuring the “size” of the knotted part in the polygon is to find a minimal length sub-walk of the polygon which has knot type  $K$ , and then define the size of the knot to be the length of this sub-walk. This of course requires a method for assigning a

knot-type to an open chain; there are various ways to do so (see ref. 7 and 16 for reviews). If  $K$  is composite, then one can use a similar approach for each of the components. We refer to this measurement of knot-size as arclength knot-size.

Motivated by the arclength knot-size definition, in this section we introduce an approach for classifying the knotted parts of polygons in  $\mathbb{T}$  as either “local” or “non-local”. We will also introduce another measure for knot-size which is particularly suited to polygons in narrow tubes. References will be made to Fig. 2 for illustration.

Let  $\pi$  be a polygon in  $\mathbb{T}$  with span  $s$ . Thus  $\pi$  is embedded in  $\mathbb{T}$  between the planes  $x = 0$  and  $x = s$ . For half-integers  $k \in \mathbb{Z} + \frac{1}{2}$  with  $0 < k < s$ , we say that  $\pi$  has a 2-section at  $x = k$  if the plane  $x = k$  intersects  $\pi$  at exactly two points. (Equivalently,  $\pi$  has exactly two edges in the  $x$ -direction between  $x = k - \frac{1}{2}$  and  $x = k + \frac{1}{2}$ .) See Fig. 2(b) where  $k = 1/2$  or  $k = 11/2$ . If  $\pi$  has  $m$  2-sections, let  $t(\pi) = (t_1, \dots, t_m)$  be the (ordered) set of  $x$ -values at which they occur. Clearly  $m \leq s$ ; if  $m = 0$  then  $t(\pi)$  is empty.

The 2-sections of a polygon  $\pi$  in  $\mathbb{T}$  naturally partition it into a sequence of “segments”. Moreover, if  $\pi$  has prime knot-type, then typically the segment of  $\pi$  which contains the “knotted part” will lie between two successive 2-sections. See for examples Fig. 2(b) and (c) for polygons with trefoil knot-type. It is this idea which will allow us to locate, measure and classify knot components within polygons in  $\mathbb{T}$ .

Take  $\pi$  in  $\mathbb{T}$  with  $m \geq 2$ . The 2-sections of  $\pi$  partition it into a sequence of segments that we call connect-sum patterns (cs-patterns for short). Then for any  $1 \leq i < m$ , the segment of  $\pi$  between  $x = t_i$  and  $x = t_{i+1}$  is called a proper cs-pattern of  $\pi$ . If  $m \geq 1$  then the segment of  $\pi$  between  $x = 0$  and  $x = t_1$  is called the start cs-pattern of  $\pi$ ; likewise, the segment of  $\pi$  between  $x = t_m$  and  $x = s(\pi)$  is called the end cs-pattern of  $\pi$ . See Fig. 2(b) and (c) for examples with  $m = 2$ . Note that proper, start and end cs-patterns are examples of, respectively, proper, left-most and right-most patterns as defined in ref. 10. Here we define the span of a (resp. start, end) cs-pattern to be  $t_{i+1} - t_i + 1$  (resp.  $t_1 + \frac{1}{2}, s(\pi) - t_m + \frac{1}{2}$ ).

Any proper cs-pattern  $\sigma$  (between  $x = t_i$  and  $x = t_{i+1}$ , for some  $i \geq 1$ ) is the union of two “strands” (self-avoiding walks)  $\sigma_1$  and  $\sigma_2$ , each extending from the left end of the pattern to the right. On the left side of the plane  $x = t_i$ , joining the two left ends of  $\sigma_1$  and  $\sigma_2$  to each other and then, on the right side of the plane  $x = t_{i+1}$ , joining their two right ends to each other, yields what we call the denominator closure of the cs-pattern. See Fig. 3(a). Note here that if the overall polygon has knot-type  $K$  and one of its proper cs-patterns has denominator closure of knot-type  $K'$ , then  $K'$  must be part of the knot-decomposition of  $K$ . Let  $DC(\sigma)$  be the knot-type of the denominator closure of  $\sigma$ . If  $DC(\sigma) \neq 0_1$  (the unknot), then we say that  $\sigma$  is a knot pattern with knot-type  $DC(\sigma)$ .

Alternatively, the two endpoints of  $\sigma_1$  (resp.  $\sigma_2$ ) can be reconnected to each other (outside of  $\mathbb{T}$ ) to form a (possibly separable) link. We call this the numerator closure of  $\sigma$ . See Fig. 3(b). Here we are not interested in the overall link-type of the numerator closure; we instead only care about the knot-types of its two components. Let  $NC_1(\sigma)$  (resp.  $NC_2(\sigma)$ ) be the knot-type of the closure of  $\sigma_1$  (resp.  $\sigma_2$ ).

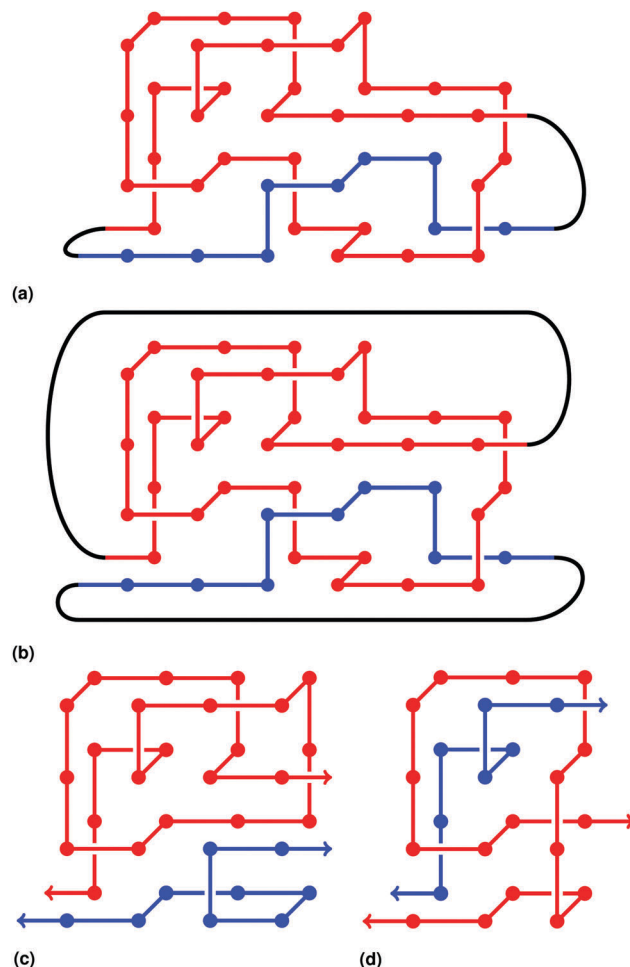


Fig. 3 (a) An illustration of the denominator closure of a proper cs-pattern  $\sigma$ . The blue strand corresponds to  $\sigma_1$  and the red to  $\sigma_2$  and their union is proper cs-pattern  $\sigma$ . The denominator closure is obtained by adding the black arcs and yields a closed curve with knot-type  $5_1$ , i.e.  $DC(\sigma) = 5_1$ . (b) The numerator closure of the same pattern. The numerator closure gives a link with one component a  $3_1$  ( $\neq 5_1$ ) knot and the other an unknot ( $0_1 \neq 5_1$ ); hence this is a non-local knot pattern. Here  $NC_1(\sigma) = 0_1$  and  $NC_2(\sigma) = 3_1$ . (c) A local trefoil knot pattern that can occur in a Hamiltonian polygon. (d) A non-local trefoil knot pattern that can occur in a Hamiltonian polygon.

We are now prepared to give our definitions of local and non-local knot patterns. Let  $\sigma$  be a cs-pattern of a polygon. If  $\sigma$  is a proper pattern with  $DC(\sigma) = K \neq 0_1$  but such that  $NC_1(\sigma) \neq K \# K'$  and  $NC_2(\sigma) \neq K \# K''$  for any  $K'$  or  $K''$  (i.e.  $K$  is not in the knot decomposition of either  $NC_1(\sigma)$  or  $NC_2(\sigma)$ ), then the knot  $K$  cannot be discovered by examining only one of the strands of  $\sigma$ . In this case both strands are needed to detect  $K$  and hence since the two strands of  $\sigma$  could potentially be far apart along the contour of the entire polygon, we define  $\sigma$  to be a non-local knot pattern. For example, the denominator closure in Fig. 3(a) is a  $5_1$  knot while the components of the numerator closure in Fig. 3(b) are a trefoil and an unknot; the corresponding proper cs-pattern is therefore a non-local knot pattern. For all other cases we classify  $\sigma$  as a local knot pattern. Note that there are examples of cs-patterns which we classify here as “local” which might be more appropriately classified as “non-local”



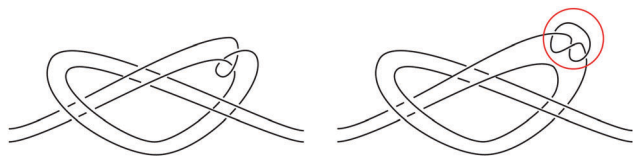


Fig. 4 Two examples of local knot patterns. The right one has a red 2-sphere intersecting the pattern in two points and surrounding a local knot.

(see Fig. 4 and the related discussion below), however, for small tube sizes such cs-patterns have a relatively low probability of occurrence.

The decomposition of polygons into cs-patterns leads to a second measure for knot-size in a polygon as follows. A polygon with prime knot-type  $K$  has at most one knot pattern (either local or non-local). If it has a knot pattern, then we define its connect-sum knot-size to be the number of edges in the knot pattern. If there is no knot pattern in the polygon, then the length of the whole polygon is used. A similar approach can be used for a composite knot  $K$ , where we use the number of edges in each (if any) knot pattern to determine a total connect-sum knot-size.

In this section we have defined a classification scheme for local knot patterns that is relatively simple to implement for polygons in tubes, however, this classification is not necessarily consistent with the topological definition of local knot. In particular, a topologically precise approach to defining a local knot pattern is as follows. For a given knot pattern, if we can find a 2-sphere that intersects the pattern in exactly two points and such that it surrounds a knotted arc, then the knot pattern contains a local knot. The right side of Fig. 4 shows an example of a knot pattern with such a 2-sphere. However, by our classification scheme, the left side of Fig. 4 shows an example of a knot pattern  $\sigma$  which we are currently classifying as a local knot pattern even though there is no 2-sphere satisfying the property just described. In particular,  $DC(\sigma) = NC_1(\sigma) = NC_2(\sigma) = 3$ , however, there is no 2-sphere that isolates (by surrounding it at the exclusion of anything else) either of these trefoils and intersects the pattern in only two points.

Because the number of edges needed to create patterns such as that shown in Fig. 4 (left) is large, such patterns are not very likely to occur and we do not believe that their existence will affect our conclusions.

In the next section we explore numerically the occurrence probabilities associated with local and non-local cs-patterns for small tube sizes and specific knot-types and we provide evidence that non-local knot patterns are more probable than local ones. However, even though a polygon  $\pi$  with knot-type  $K \neq 0_1$  contains a knot pattern of a given type (local or non-local), this does not necessarily guarantee that the polygon is locally or non-locally knotted (using the standard arclength classification); this aspect is explored more thoroughly and theoretically in Section 5.

## 4 Exact and simulation results

Exact generation was used to determine all smallest-span non-local and local trefoil knot patterns in the  $2 \times 1$  and  $3 \times 1$  tubes.

Table 1 Numbers of trefoil patterns of smallest spans in the  $2 \times 1$  and  $3 \times 1$  tubes, for all and Hamiltonian (Ham.) polygons

Tube size	Span	Non-local	Local	Ham. non-local	Ham. local
$2 \times 1$	6	116	0	32	0
	7	5584	304	668	80
	8	141 292	14 932	8020	1388
$3 \times 1$	4	1964	0	232	0
	5	762 984	29 272	17 568	1448

(See Fig. 2 for examples of such patterns in the  $2 \times 1$  tube and Fig. 3 in the  $3 \times 1$  tube.) Counts are shown in Table 1. In all cases the number of non-local knot patterns greatly exceeds the number of local knot patterns, suggesting that non-local trefoil knot patterns may be more likely to occur than local ones. To explore whether this conclusion depends on the model used (fixed-edge or fixed-span), limiting probabilities of occurrence of each type of pattern were determined under each of the distributions  $\mathbb{P}_n^{(ed,f)} (-\infty < f < \infty, n \rightarrow \infty)$  and  $\mathbb{P}_s^{(sp,g)} (-\infty < g < \infty, s \rightarrow \infty)$ . These limiting probabilities can be determined (see for example Eng<sup>14</sup>) from the eigenvalues and eigenvectors of the transfer-matrix. Fig. 5 shows the results for the  $3 \times 1$  tube. In this figure, for the fixed-edge model ( $\mathbb{P}_n^{(ed,f)}$ ),  $\mathbb{P}_{3_1}^{ed,L}(f)$  denotes the limiting ( $n \rightarrow \infty$ ) probability of occurrence of a smallest local trefoil knot pattern at a section of a polygon and  $\mathbb{P}_{3_1}^{ed,NL}(f)$  denotes the corresponding probability for the non-local patterns. Similarly, for the fixed-span model ( $\mathbb{P}_s^{(sp,g)}$ ),  $\mathbb{P}_{3_1}^{sp,L}(g)$  denotes the limiting ( $s \rightarrow \infty$ ) probability of occurrence of a smallest local trefoil knot pattern at a section of a polygon and  $\mathbb{P}_{3_1}^{sp,NL}(g)$  denotes the corresponding probability for the non-local patterns. Further note that Fig. 5 shows the results for the fixed-edge probabilities with the horizontal axis corresponding to  $f$  while for the fixed-span probabilities it corresponds to  $-g$ . The latter was done to make an easier comparison between the models, since positive values of  $f$  and negative values of  $g$  both have a stretching effect on polygons. Although not shown here, the observed trends were similar for the  $2 \times 1$  tube. The results are summarised below.

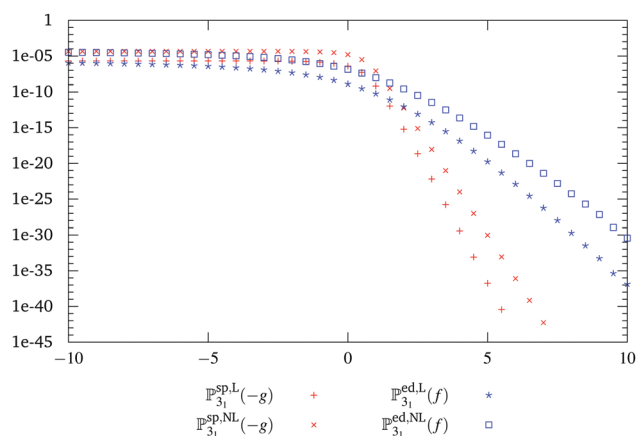


Fig. 5 log scale plot of the probabilities of the smallest local and non-local trefoil patterns in the  $3 \times 1$  tube, as functions of  $f$  (blue) and  $-g$  (red).

**Result 1.** For  $\mathbb{T}_{2,1}$  and  $\mathbb{T}_{3,1}$ , the limiting occurrence probability of the smallest-span non-local trefoil knot patterns is greater than that for the smallest-span local trefoil knot patterns, regardless of the strength of the force  $f$  or the edge density weight  $g$ .

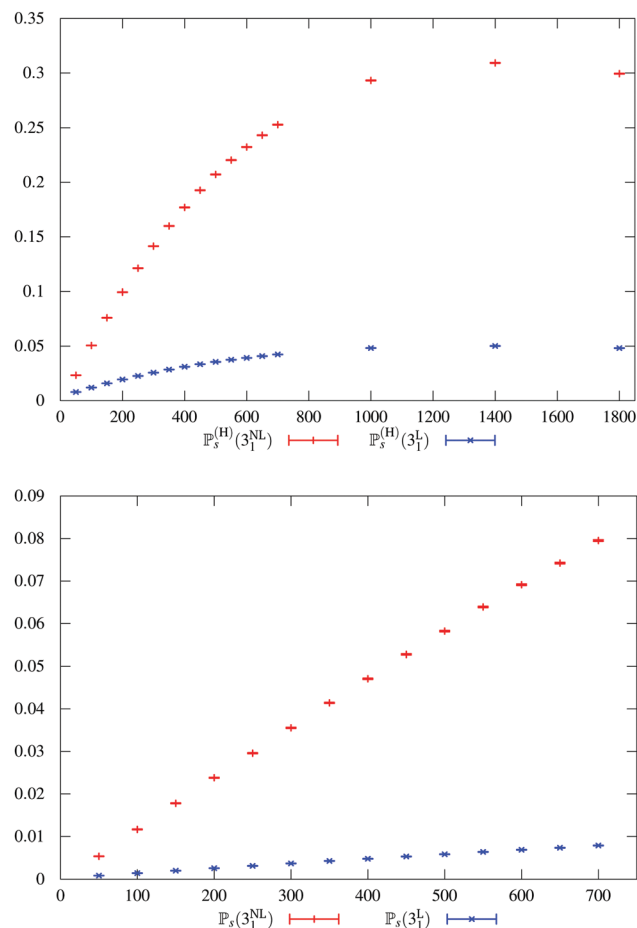
Determining the knot-type of a polygon or a knot pattern requires the whole polygon or knot pattern. However, since the numbers of polygons and knot patterns in a tube grow exponentially with either span or the number of edges, we limited the exact generation analysis to the case of the smallest trefoil patterns. To explore further whether the trend observed for the smallest trefoil patterns holds for other knots, a Monte Carlo approach was developed to generate random polygons in the tube, based on a method of Alm and Janson.<sup>17</sup> The Monte Carlo method is also based on transfer-matrices and can be used to generate a set of independent and identically distributed polygons from any of the distributions  $\{\mathbb{P}_n^{(\text{ed},f)}, \mathbb{P}_s^{(\text{sp},g)}\}$  provided that the transfer-matrix associated with  $G_T(f,g)$  is known. Details of the approach will be published elsewhere.

Based on the exact results of Fig. 5, we focused on the fixed-span model at  $g = 0$  where the probabilities of the smallest trefoil patterns were greatest (compared to the other models). Similarly we focus on the  $3 \times 1$  tube, since knots are far more common than in  $2 \times 1$  while the transfer matrices are small enough as to make simulations and enumerations reasonably efficient.

For the Monte Carlo results, we begin by investigating the probabilities of some simple knots. In Fig. 6 we plot the probabilities of local and non-local trefoils in the fixed-span ensemble, for both Hamiltonian and all polygons. Specifically we plot  $\mathbb{P}_s(K^{\text{NL}})$  and  $\mathbb{P}_s(K^{\text{L}})$  which are respectively the observed proportions of span  $s$  polygons which have knot-type  $K = 3_1$  and contain a non-local (NL) or local (L) knot pattern. We also plot the corresponding observed proportions of span  $s$  Hamiltonian polygons:  $\mathbb{P}_s^{(\text{H})}(K^{\text{NL}})$ ,  $\mathbb{P}_s^{(\text{H})}(K^{\text{L}})$ . The corresponding data for figure-eight knots ( $K = 4_1$ ) is illustrated in Fig. 7. The relative frequencies of non-local knots for trefoils ( $3_1$ ), figure-eight knots ( $4_1$ ),  $5_1$  and  $5_2$  knots are illustrated in Fig. 8. For example, the relative frequency of non-local trefoil knots amongst span  $s$  trefoil polygons is  $\mathbb{P}_s(3_1^{\text{NL}}|3_1) = \mathbb{P}_s(3_1^{\text{NL}})/\mathbb{P}_s(3_1)$ . Although not shown, similar trends were observed for the  $2 \times 1$  tube. Our observations lead to the following conclusion.

**Result 2.** For  $\mathbb{T}_{2,1}$  and  $\mathbb{T}_{3,1}$ , based on i.i.d. samples from the limiting distribution of  $\mathbb{P}_s^{(\text{sp},0)}$  over a range of spans  $s$  (10 million polygons per span), we observe that the probability of occurrence of a non-local  $3_1$ ,  $4_1$ ,  $5_1$  or  $5_2$  knot pattern is significantly greater than that of the corresponding local knot pattern (e.g. Fig. 6 (bottom) and Fig. 7 (bottom)). The same holds for Hamiltonian polygons sampled from the limiting distribution of  $\mathbb{P}_s^{(\text{H})}$  (e.g. Fig. 6 (top) and Fig. 7 (top)). Furthermore, for sufficiently long polygons in  $\mathbb{T}_{3,1}$ , for each of these knot-types the proportion of non-local patterns amongst all observed knot patterns of that type is greater than 85% (see Fig. 8).

Next, we consider the average span of knot patterns within polygons. Recall that we define the span of a knot pattern to include the two 2-sections which bound it on the left and right. In Fig. 9 we plot the average span of non-local and local trefoil



**Fig. 6** Plots of the probabilities of non-local (red) and local (blue) trefoil knots, for Hamiltonian (above) and all (below) polygons in the  $3 \times 1$  tube, sampled uniformly from the fixed-span ensemble ( $g = 0$ ). The horizontal axis is span  $s$ . The (barely visible) error bars represent 95% confidence intervals.

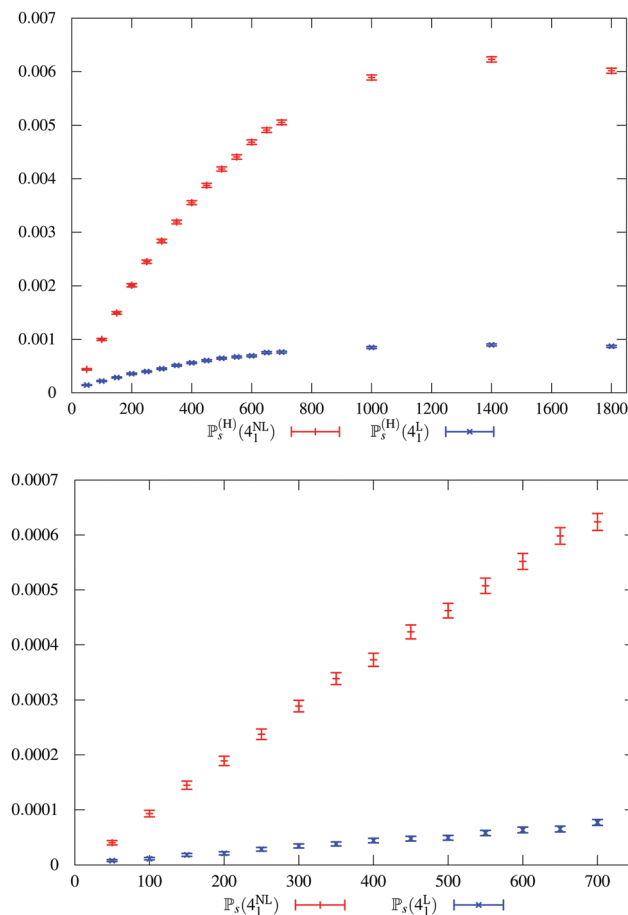
and figure-eight knot patterns, for both Hamiltonian and all polygons sampled uniformly from the fixed-span ensemble. Although not shown, similar trends were observed for the  $2 \times 1$  tube. Our observations lead to the following conclusion.

**Result 3.** For  $\mathbb{T}_{2,1}$  and  $\mathbb{T}_{3,1}$ , based on i.i.d. samples from the limiting distribution of  $\mathbb{P}_s^{(\text{sp},0)}$  over a range of spans  $s$  (10 million polygons per span), we observe that the average spans of non-local  $3_1$ ,  $4_1$ ,  $5_1$ , and  $5_2$  knot patterns are smaller than those of the corresponding local knot patterns (e.g. Fig. 9 (bottom)). The same holds for Hamiltonian polygons sampled from the limiting distribution of  $\mathbb{P}_s^{(\text{H})}$  (e.g. Fig. 9 (top)).

The results of this section provide strong numerical evidence that at least for small tube sizes, non-local knot patterns are more likely and have on average shorter span than local knot patterns. In the next section we discuss to what extent results related to this can be proved.

## 5 Theoretical results

In this section we state results that can be proved related to the occurrence of non-local and local knot patterns in tubes.



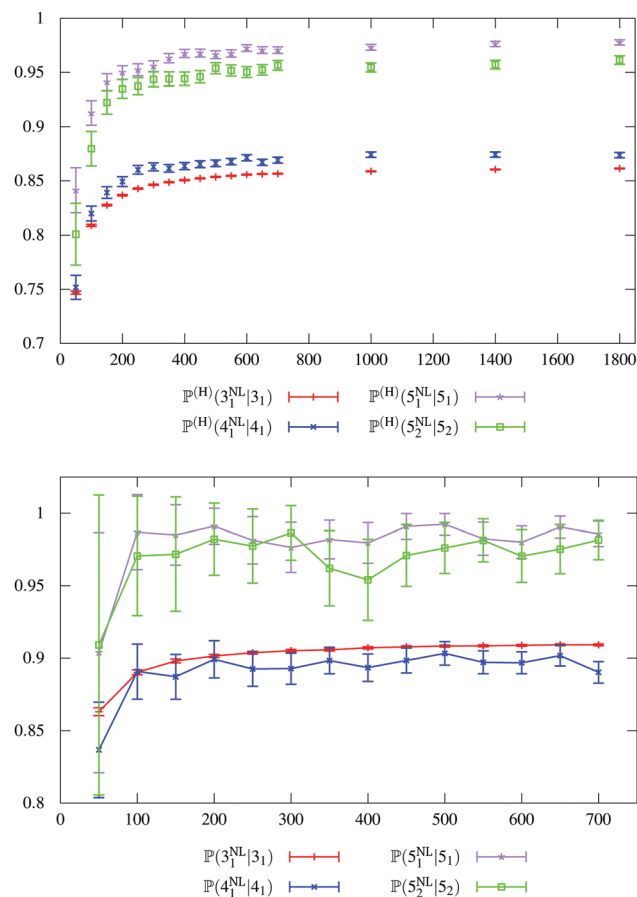
**Fig. 7** Plots of the probabilities of non-local (red) and local (blue) figure-eight knots, for Hamiltonian (above) and all (below) polygons in the  $3 \times 1$  tube, sampled uniformly from the fixed-span ensemble ( $g = 0$ ). The horizontal axis is span  $s$ . Error bars represent 95% confidence intervals.

We focus on the statement of the results and leave most details of any proofs to the ESI.<sup>†</sup> First we present results related to which knot patterns can occur in a tube of a given size and what is known about the minimum span of such knot patterns. Then we present results on the probability of occurrence of knot patterns.

Because of the strict geometric confinement, not all knots are embeddable in a given  $\mathbb{T}$ ; the dimensions of the tube determine whether a particular knot is embeddable.<sup>18</sup> We present here a general statement; a more precise (but technical) version is stated and proved in the ESI.<sup>†</sup>

**Result 4.** For any given knot type  $K$ ,  $K$  admits a proper non-local knot pattern in a tube  $\mathbb{T}_{L,M}$  for  $L, M$  sufficiently large, and admits a proper local knot pattern in a tube  $\mathbb{T}_{L',M'}$  for  $L', M'$  sufficiently large. Any tube  $\mathbb{T}$  which accommodates a local knot pattern for  $K$  also accommodates a non-local knot pattern.

Note that this result leaves open the possibility that for a given knot  $K$  with a non-local knot pattern in  $\mathbb{T}_{L,M}$  there might not be an associated local knot pattern in  $\mathbb{T}_{L,M}$ . For the case of a  $2 \times 1$  tube, however, we establish a more definitive result – namely, if the knot is embeddable in the tube then there also exist both non-local and local knot patterns for the knot in the



**Fig. 8** Relative frequencies of non-local knots for trefoils (red), figure-eight knots (blue),  $5_1$  knots (purple) and  $5_2$  knots (green), for Hamiltonian (above) and all (below) polygons in the  $3 \times 1$  tube, sampled uniformly from the fixed-span ensemble. Note that the lines joining the points in the second plot have been added only to aid the reader, and do not indicate any additional data.

tube. Fig. 2(a) shows a trefoil polygon in a  $2 \times 1$  tube that yields a non-local knot pattern (see Fig. 2(b)). Fig. 2(c) shows a local knot pattern in the same tube; note that the span of this local knot pattern is one greater than that shown in Fig. 2(b). For the  $2 \times 1$  tube, the arguments used by Ishihara *et al.*<sup>18</sup> can also be extended to prove that this difference in span holds for the smallest knot patterns of knots with up to 5 crossings. In summary the following result can be proved.

**Result 5.** Given a prime knot  $K \neq 0_1$  that can occur in a  $2 \times 1$  tube, there exists at least one proper local knot pattern and at least one proper non-local knot pattern. Furthermore, at least for  $K \in \{3_1, 4_1, 5_1, 5_2\}$ , the span of a smallest proper local knot pattern of  $K$  in  $\mathbb{T}_{2,1}$  is greater than that of a smallest proper non-local knot pattern of  $K$  in  $\mathbb{T}_{2,1}$ .

For the  $3 \times 1$  tube any knot that can occur in  $\mathbb{T}_{2,1}$  will also have non-local and local knot patterns; see Fig. 3 for some examples. We note that for the  $3 \times 1$  tube, by an exhaustive search, we have determined that the span of a smallest local knot pattern of  $3_1$  is also one greater than that of a smallest proper non-local knot pattern of  $3_1$  (Fig. 3(c) and (d) are examples of such smallest knot patterns for Hamiltonian

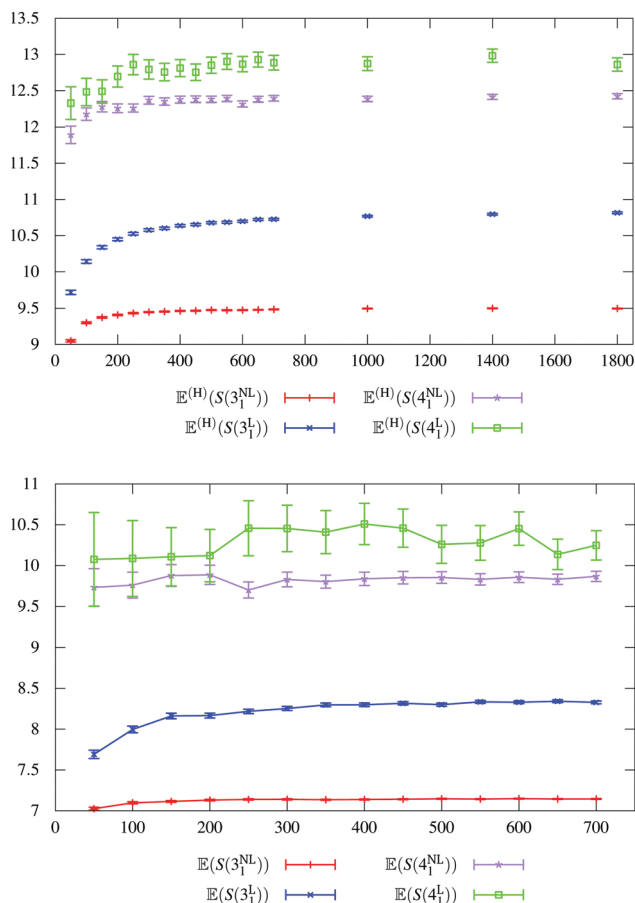


Fig. 9 Average span of non-local trefoil (red), local trefoil (blue), non-local figure-eight (purple) and local figure-eight (green) knot components versus overall polygon span, for Hamiltonian (above) and all (below) polygons in the  $3 \times 1$  tube, sampled uniformly from the fixed-span ensemble.

polygons in the  $3 \times 1$  tube). We also looked at the smallest prime knots which cannot occur in  $\mathbb{T}_{2,1}$  but can in  $\mathbb{T}_{3,1}$ :  $8_5, 8_{10}, 8_{15} - 8_{21}$ . For all these knots, it is possible to construct a local pattern.

The fact that the span of smallest non-local knot patterns is smaller than that of local knot patterns provides a partial explanation for Result 3, that the average span of a non-local knot pattern is smaller than that of a corresponding local knot pattern for small tube sizes. However, for a large enough tube size, the span of smallest non-local and local knot patterns are expected to be the same; see Fig. 10 which shows, on the left, a shortest trefoil arc which is part of a span-3 local trefoil knot pattern within  $\mathbb{T}_{3,3}$ , and shows on the right a smallest non-local trefoil knot pattern in  $\mathbb{T}_{3,3}$  also having span 3. Note, however, for this case the number of edges in the local knot pattern is greater than that for the non-local pattern; this may lead to non-local trefoil patterns being more likely to occur than local ones even in larger tube sizes, *i.e.* that Result 2 could continue to hold.

We next discuss some results about the likelihood of each type of pattern.

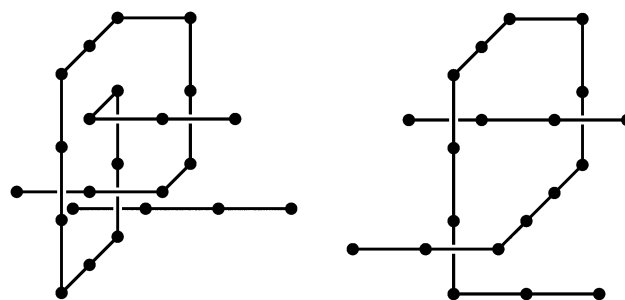


Fig. 10 The shortest local (left) and non-local (right) knot patterns. Both fit in  $\mathbb{T}_{3,3}$ .

There are known “pattern theorems” available for both the fixed-edge and fixed-span models studied here (see ref. 11 and 15), as well as for Hamiltonian polygons (see ref. 14). The theorems focus on proper polygon patterns (see ref. 10 for more precise definitions) which include the proper knot patterns defined here. Given a model and a proper pattern  $P$  which can occur in a polygon of the model in  $\mathbb{T}$ , a pattern theorem establishes that there exists an  $\varepsilon_P > 0$  such that, for  $n$  sufficiently large, all but exponentially few  $n$ -edge polygons contain more than  $\varepsilon_P n$  copies of  $P$ . From such theorems it is known that the knot-complexity of polygons grows as polygon “size” grows (size could be measured in terms of edges or span), so that a typical polygon will have a highly-composite knot-type  $K = K_1 \# K_2 \# \dots \# K_r$ .<sup>19</sup> Different prime components of the knot could occur as knot patterns in the polygon in a variety of ways. Our interest here is to investigate how often they are occurring as “local” knots *versus* non-locally.

To define “local” knotting requires the definition of a knot-size measure (we have given two possible measures: arclength and connect-sum knot-size) but also a comparison of knot-size to polygon size. Here we consider that a polygon’s size  $m$  is growing without bound and say that it is non-locally knotted with respect to arclength knot-size if at least one of the knots  $K_i$  in its prime knot decomposition has arclength knot-size  $a_{K_i} = O(m)$ . In this case we say  $K_i$  occurs non-locally in the polygon or is non-locally knotted, and otherwise we say that  $K_i$  occurs locally or is locally knotted. Thus a polygon can be both non-locally and locally knotted (with respect to arclength knot-size) depending on the occurrence of each of its prime components. Corresponding definitions can apply to the case of the connect-sum knot-size. However, to distinguish this case, we say a polygon is “loosely” knotted (or contains a loose knot) if at least one of the knots  $K_i$  in its prime knot decomposition has connect-sum size  $b_{K_i} = O(m)$ ; otherwise, we say  $K_i$  occurs as a tight knot or tightly. See Fig. 11.

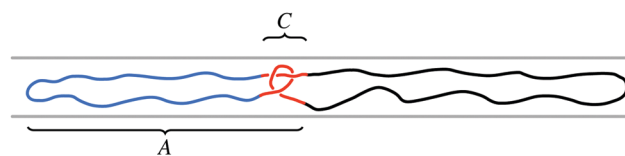


Fig. 11 A schematic of the non-local trefoil pattern from Fig. 10 in a polygon. This knot is tight (because the connect-sum size  $C$  is small) but non-local (because the arclength size  $A$  is large).



To explain these definitions more clearly, consider the non-local trefoil pattern  $P$  of Fig. 2(b). If this pattern occurs in a trefoil polygon  $\pi$ , then  $\pi$  must be formed from a connect sum of unknotted polygons with  $P$ . Suppose the size of  $\pi \gg$  the span of  $P$ . If  $P$  occurs near the left end (or right end) of  $\pi$ , then the arclength knot-size will be short compared to the size of  $\pi$  and even though we have classified  $P$  as a non-local knot pattern, we will say that  $P$  has occurred in a “local” way in  $\pi$  and that  $\pi$  is locally knotted. If instead  $P$  occurs in the “middle” of  $\pi$  (*i.e.* half-way along the span) then we will say that  $P$  has occurred in a “non-local” way in  $\pi$  and that  $\pi$  is non-locally knotted (since the arclength size of the knot is proportional to the size of  $\pi$ ). In contrast, because the connect-sum knot-size of  $P$  is small compared to the size of  $\pi$ , no matter where it occurs in  $\pi$  we will say that  $3_1$  occurs as a tight knot. On the other hand, for the local trefoil pattern of Fig. 2(c), no matter where it occurs in a very large sized trefoil polygon  $\pi$ , it will always have a small arclength knot-size as well as a small connect-sum knot-size and hence  $\pi$  will be considered to be both locally and tightly knotted.

Since the pattern theorems hold for proper knot patterns, they tell us that for each of the models in question, any proper knot pattern which can occur in a polygon (or Hamiltonian polygon) will occur with a positive density as polygon size grows (where, depending on the model, size is measured by edges or by span). In particular for any  $\mathbb{T}$ , polygons which contain no knot patterns are exponentially rare. Furthermore we have, by considering the local trefoil knot pattern of Fig. 2(c) (or an inflated version for the Hamiltonian polygon cases), the following result.

**Result 6.** All but exponentially few sufficiently long polygons in  $\mathbb{T}$  with  $M \geq L \geq 2$  or  $M \geq 2, L = 1$  for any of the three models (fixed-length, fixed-span and Hamiltonian) are both locally and tightly knotted.

Note that this does not preclude the same polygons from being non-locally or loosely knotted – we only know that the knot-types of the polygons are highly complex and that some of the knots in the knot decomposition will be local trefoil knot patterns as in Fig. 2(c).

Applying the above argument to proper non-local knot patterns also leads to the following.

**Result 7.** All but exponentially few sufficiently long polygons in  $\mathbb{T}$  with  $M \geq L \geq 2$  or  $M \geq 2, L = 1$  are non-locally knotted.

Meanwhile for self-avoiding walks in  $\mathbb{T}$ , at least for  $f = 0$ , the scenario depicted in Fig. 1(e), in which a non-local knot pattern occurs in a walk in a non-local way, is exponentially rare. To see this, first note that any polygon can be turned into a walk by removing one edge. Thus from (1) for  $f = 0$ , we know that polygons are exponentially rare in the set of walks. Next consider the subset of  $n$ -step walks in  $\mathbb{T}$  which contain a non-local knot-type  $K$  pattern at a location in the walk such that the arclength of  $K$  is  $\alpha n$  for some  $\alpha > 0$  (as in Fig. 1(e)). Each such walk can be decomposed into a polygon with at least  $\alpha n$  edges (and having  $K$  in its knot-decomposition) and a walk with length at most  $(1 - \alpha)n$ . Thus this subset of walks will have an exponential growth rate which is strictly less than all walks in  $\mathbb{T}$ . This gives the following result.

**Result 8.** Self-avoiding walks which contain a non-local knot pattern in a non-local way are exponentially rare in the set of all walks in  $\mathbb{T}$ .

(Note that this result does not contradict the pattern theorem for walks in  $\mathbb{T}$  proved by Soteros and Whittington<sup>13</sup> because the proper knot patterns defined here are not examples of proper walk patterns.)

## 6 Practical implications and DNA models

In this section we will briefly discuss the connections between this work and other models of confined DNA. Before proceeding, we again reiterate that for simplicity our model does not account for any longer-range electrostatic interactions.

### 6.1 Tube dimensions

Self-avoiding polygons on  $\mathbb{Z}^3$  have a natural ring structure, but when confined to a tube of diameter much smaller than the polygon length, they become essentially linear objects. It is thus natural to look at how correlations decay along the length of the tube; and since our Monte Carlo sampling method constructs polygons one cs-pattern at a time using a transfer matrix, it is also straightforward to study this.

For the fixed-span model, we have measured the number of edges in a section as well as the centre of mass of those edges, and looked at how these quantities correlate in two sections separated by distance  $s$  (within a long polygon). The centre of mass has slightly stronger correlations, but still very weak: the correlation between two adjacent sections is  $\rho \approx -0.061$ , and this decays exponentially for  $s > 1$ . By adapting a formula for the autocorrelation time of a stationary Markov chain,<sup>20</sup> we obtain an estimate for the “correlation span” along the length of the tube:  $s_c \approx 1.12$ . For Hamiltonian polygons, a similar analysis leads to  $s_c \approx 1.66$ . Now the average density of a polygon in the  $3 \times 1$  tube is about 2.05 edges per strand per unit span,<sup>14</sup> while for Hamiltonian polygons it is exactly 4. We thus get the rough “correlation lengths”  $\ell_c \approx 2.31$  and  $\ell_c \approx 6.63$  respectively.

We take the “diameter” of the  $3 \times 1$  tube to be  $D = \sqrt{3} = 1.732$ . Since double-stranded DNA has a persistence length of about 50 nm,<sup>21</sup> we arrive at the estimate that  $D$  roughly corresponds to the range 10–40 nm, so that  $D/\ell_c < 1$ . In Micheletti and Orlandini<sup>3</sup> evidence is provided that, at least for high salt solutions, there is a crossover from the Odijk scaling region (where  $D/\ell_c < 1$ ) to the de Gennes scaling regime at channel widths of about 85 nm. Our model results for small tube sizes are consistent with a scaling regime that is well below the de Gennes scaling regime.

### 6.2 Other models

Micheletti and Orlandini<sup>3</sup> model DNA in a nanochannel using a semi-flexible chain of cylinders (with parameters chosen to match those of double stranded DNA in concentrated solutions of monovalent salts) and confine the chain to cylindrical channels of various widths. They present numerical results

on knotting probabilities as a function of chain length and channel width. They also explore the size of knots using an arclength measure; Fig. 4(c)–(f) in that paper show results for the average lengths and the distribution of lengths of trefoil knots in linear and circular chains as functions of total chain length as well as the effective nanochannel diameter. Considering Fig. 4(c) from that paper and focussing on the smallest effective diameter, it can be seen that for linear DNA the average trefoil arclength is relatively independent of total chain length for sufficiently large linear DNA; this is consistent with trefoils occurring as local knots (as predicted in our Result 8). In contrast, considering Fig. 4(d) from that paper and again focussing on the smallest effective diameter, it can be seen that for circular DNA the average trefoil arclength is on the order of 1/4 total chain length; this is consistent with the trefoil occurring as a small non-local knot pattern randomly distributed through the chain (consistent with our result that non-local knot patterns are more likely than local knot patterns in polygons). To make better comparisons between these two models, however, it would be necessary to include more DNA-like interactions into our lattice model as well as to explore larger tube sizes.

Our definitions of local and non-local knot patterns correspond very closely with a recent numerical study by Suma and Micheletti.<sup>6</sup> In that work, the translocation of knotted 10 kbp DNA rings through a nanopore of diameter 10 nm was simulated. The authors found two different modes of knot translocation, which they called single- and double-filament (see Fig. 4(C) and (D)).<sup>6</sup> These correspond to our definitions of local and non-local knot patterns respectively. For that (non-equilibrium) model, the relative frequencies of these modes (measured as the knot passes through the pore) depended on the length of the DNA molecule and the initial conditions prior to translocation. However, it was observed that the translocation time for both single- and double-filament knots was about the same, and thus this cannot be used as a method for distinguishing the two modes. There are many differences between the model presented here and that in Suma and Micheletti,<sup>6</sup> and we do not attempt to make a quantitative comparison. We expect, however, that the mathematical framework developed here for classifying these two types of knot modes will prove useful for further analysis of these models.

### 6.3 DNA experiments

Our results are also connected to recent experimental work,<sup>5</sup> which demonstrated that knots in DNA can be detected by passing the molecule through a nanopore and observing when and for how long the pore is obstructed. While both local and non-local knot patterns can be detected, as observed in ref. 6, it does not seem possible to distinguish the two modes of knotting using this method, as the obstruction events take about the same time to occur. On the other hand, if it were possible to measure the arclength of the knot as well as the translocation time, then our results would suggest that those knots with an arclength which is inconsistently long, in comparison to

the translocation time, could be classified as occurring in the double-filament mode (*i.e.* occurring as a non-local knot pattern).

The biological relevance of these results is presently unclear. The nature of a knot in a DNA molecule is biologically important, as it can have an effect on DNA replication or the effectiveness of enzymes like topoisomerase in untangling the knot. Whether local or non-local knot patterns can be more easily untangled is an open question, and one we intend to study further.

We also mention some experimental work on stretched DNA knots.<sup>22,23</sup> In those cases the authors studied linear DNA, and observed knots in stretched molecules by microscopy techniques. It is unclear whether or not, or how, one could distinguish between the two knotting modes in any corresponding experiments for circular DNA.

### 6.4 Linking in polymers

In this work we considered only a single polygon or walk confined to a lattice tube, as a model of a single confined polymer. The topological properties of multiple molecules, like linear or circular DNA molecules or melts of rings, are also important in biological and physical systems. Methods for detecting and quantifying linking of long polymer chains have been developed, in both free solutions<sup>24</sup> and in confinement.<sup>25</sup>

The methods employed here for detecting and classifying knots in tubes can be adapted to links in tubes in a straightforward way. The “linked portion” of a pair of linked polygons can be found by finding the 2-sections of the link and examining the portions which lie between them; connecting up the loose ends in different ways (similarly to Fig. 3) is expected to enable the classification of different types of links. However, more challenging is the case where each polygon in a linked pair spans the same portion of the tube (*i.e.* there are no 2-sections in the link). This latter model has been studied theoretically in Atapour *et al.*<sup>26</sup> and is being explored now by us numerically using transfer-matrix and Monte Carlo methods to determine the linking distribution as well as to explore the open question of classifying link patterns.

## 7 Conclusion

We have used self-avoiding polygons to model ring polymers confined to narrow tubes. For this model we have used a standard approach for measuring the size of a knot to define “local” and “non-local” knotting.

We have then provided both theoretical and numerical evidence that when ring polymers are confined to very narrow tubes, at equilibrium and assuming all states are accessible, non-local knotting is more likely than local knotting (Results 1–2). This may be a consequence of the fact that non-local knot configurations, at least for the simplest knots, are on average smaller than their local counterparts (Result 3). In small tube sizes they are smaller both in span and in edge count (Result 5).

These results can be compared and contrasted with recent numerical models of DNA knots translocating through a nanopore,<sup>6</sup> and related experimental techniques for detecting

knots.<sup>5</sup> Such techniques do not appear to be sensitive enough to distinguish between these two different modes of knotting. However, being aware that two different modes are possible and determining whether one is more probable or on average tighter than the other, could lead to improved analysis of both model and experimental results.

We also provided theoretical evidence that for linear chains, non-local knotting is exponentially rare, due to the entropic disadvantage of a long bend (Result 8). This is comparable to what has been observed experimentally.<sup>5</sup>

## Conflicts of interest

There are no conflicts to declare.

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