

## Supplementary information for “Characterising knotting properties of polymers in nanochannels”

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*Note: Some Results in this document are duplicates of those in the main article – in these cases the numbering is the same. Citations and references to Figures in this document are self-contained.*

### Outlines of proofs of Results 4 and 5

**Result 4.** *For any given knot type  $K$ ,  $K$  admits a proper non-local knot pattern in a tube  $\mathbb{T}_{L,M}$  for  $L, M$  sufficiently large, and admits a proper local knot pattern in a tube  $\mathbb{T}_{L',M'}$  for  $L', M'$  sufficiently large. Any tube  $\mathbb{T}$  which accommodates a local knot pattern for  $K$  also accommodates a non-local knot pattern.*

We will prove here a more precise version of Result 4. First, we need a new definition. The *trunk* of a knot or link  $K$  is an invariant defined by  $\text{trunk}(K) = \min_E \max_{t \in \mathbb{R}} |h^{-1}(t) \cap E|$ , where  $E$  is any embedding of  $K$  in  $\mathbb{R}^3$  and  $h : \mathbb{R}^3 \rightarrow \mathbb{R}$  is any given height function [1, 3]. For another invariant, the bridge number  $b(K)$  of  $K$ ,  $\text{trunk}(K)$  satisfies  $\text{trunk}(K) \leq 2b(K)$ .

**Result 4\*.** (A) *A knot  $K$  admits a proper non-local knot pattern in  $\mathbb{T}_{L,M}$  if and only if  $\text{trunk}(K) < (L + 1)(M + 1)$ .* (B) *A knot  $K$  admits a proper local knot pattern in  $\mathbb{T}_{L,M}$  if  $\text{trunk}(K) < (L + 1)(M + 1) - 2$ .*

Given a polygon  $\pi \in \mathcal{P}_{\mathbb{T}}$ , a *hinge*  $H_k$  of  $\pi$  is the set of edges and vertices lying in the intersection of  $\pi$  and the  $y$ - $z$  plane defined by  $\{(x, y, z) : x = k\}$ . See Figure 1(a) for an example.

*Proof of Result 4\*.* (A) By [1, Theorem 1], we can construct a polygon of knot type  $K$  in  $\mathbb{T}_{L,M}$  if and only if  $\text{trunk}(K) < (L + 1)(M + 1)$ . Then we can obtain a proper knot pattern from such a polygon by opening its ends, i.e. by removing an edge or edges (as appropriate) in each of the left-most and right-most hinges. See Figure 2(a). We will show that there is a polygon which can be opened at each end to yield a proper non-local knot pattern.

First we consider the case where  $\text{trunk}(K) \geq 6$ . Take a height function  $h$  and an embedding of  $K$ ,  $\pi_K$ , in  $\mathbb{T}$  such that  $\text{trunk}(K)$  is attained and such that  $\pi_K$  has the minimal number of critical points with respect to  $h$ . We can choose one maximal point  $p$  and one minimal point  $q$  to make a proper knot pattern so that each of the two arcs of  $\pi_K - \{p, q\}$  has at least two critical points. See Figure 2(b). Let  $K_1$  and  $K_2$  be the components of the link obtained by taking the numerator closure of  $\pi_K - \{p, q\}$ . Then neither of  $K_1$  nor  $K_2$  is  $K$  by the minimality of the number of critical points of  $\pi_K$ . It follows that the pattern is non-local. We can construct a polygonal model of  $K$  satisfying the above conditions in a given tube.

Suppose  $\text{trunk}(K) = 4$ . First we consider the case where  $K$  is prime, i.e.,  $K$  is a 2-bridge knot. Take a Conway’s normal form with the minimal crossing number. Then there are at least two strings of the 4-braid corresponding to the Conway’s normal form that contain crossings. Then we can make a proper knot pattern so that both  $K_1$  and  $K_2$ , the components of the numerator closure, contain one each of such strings. Then the crossing numbers of  $K_1$  and  $K_2$  are strictly

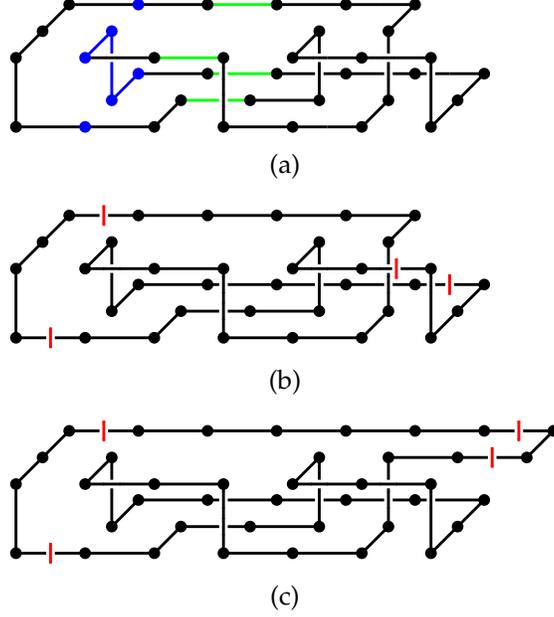


Figure 1: (a) A 36-edge polygon  $\pi$  that fits inside  $\mathbb{T}_{L,M}$  with  $L \geq 2$  and  $M \geq 1$ ; the tube extends without bound to the right and the span  $s(\pi) = 6$ . Blue vertices and edges denote the hinge  $H_1$  of  $\pi$ , and green edges denote the section  $S_3$  of  $\pi$ . (b) The locations of the two pairs of vertical red lines indicate the locations of the two 2-sections in this polygon; in this example, the polygon can be decomposed into a start unknot pattern, a proper trefoil knot pattern, and an end unknot pattern. The proper knot pattern is classified as non-local in this case. (c) A local proper knot pattern in the same tube with span 7.

less than that of  $K$ . Hence neither of  $K_1$  nor  $K_2$  is  $K$ . We can construct a polygonal model of  $K$  satisfying the above conditions in a given tube and it gives a non-local pattern. Suppose  $K$  is a composite knot. Let  $L_1$  and  $L_2$  be knots such that  $K = L_1 \# L_2$  and  $L_1$  is a prime knot. Then by the above argument, we can create a non-local pattern for  $L_1$ . By a connected sum operation, we can then construct a polygon of  $K$  that gives a non-local proper knot pattern.

(B) Suppose  $\text{trunk}(K) < (L + 1)(M + 1) - 2$ . Then by using a method of [1, Theorem 1], we can construct a polygon inside a region in  $\mathbb{T}_{L,M}$  as in Figure 2(c) (left). Then by pulling out a part as in Figure 2(c) (right) we have a local proper knot pattern.  $\square$

**Result 5.** *Given a prime knot  $K \neq 0_1$  that can occur in a  $2 \times 1$  tube, there exists at least one proper local knot pattern and at least one proper non-local knot pattern. Furthermore, at least for  $K \in \{3_1, 4_1, 5_1, 5_2\}$ , the span of a smallest proper local knot pattern of  $K$  in  $\mathbb{T}_{2,1}$  is greater than that of a smallest proper non-local knot pattern of  $K$  in  $\mathbb{T}_{2,1}$ .*

*Proof.* Any prime knot that can occur in a  $2 \times 1$  tube is 2-bridge [1]. It is well known that any 2-bridge knot is represented by Conway's normal form  $C(a_1, \dots, a_n)$ , which is a closure of a 4-braid using only two generators  $\sigma_1$  and  $\sigma_2$  [2]. Since there is no  $\sigma_3$  and the fourth string in

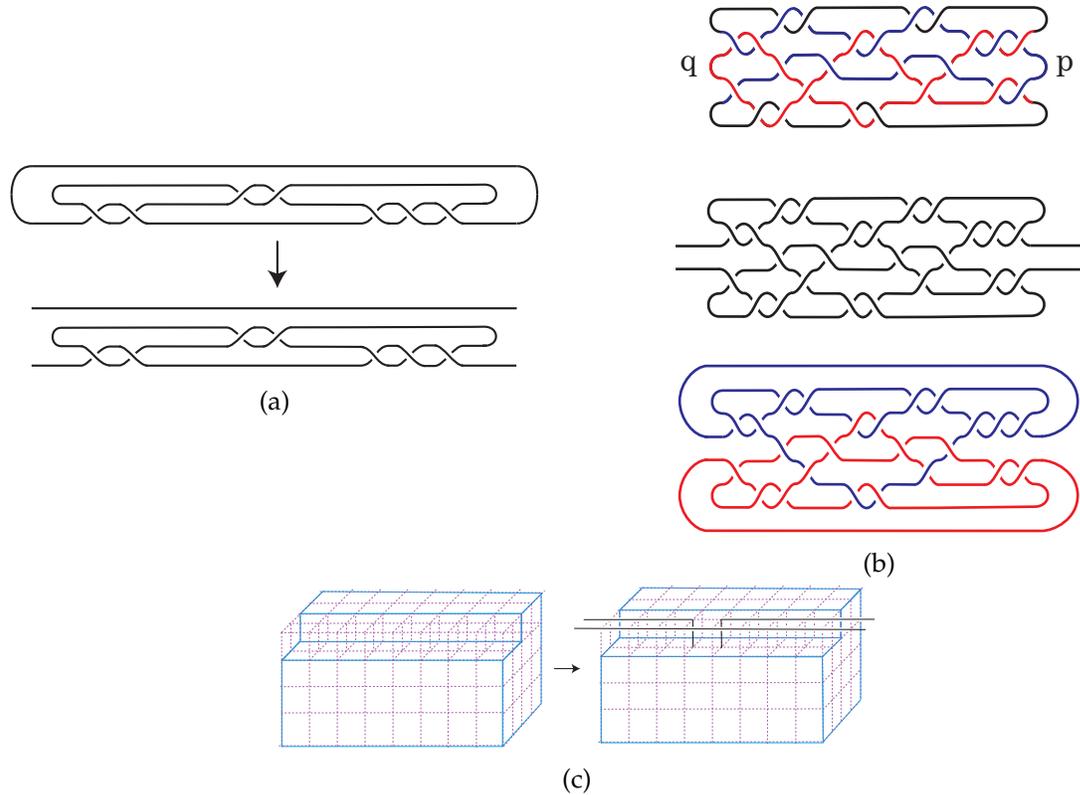


Figure 2: (a) A local  $7_5$  pattern obtained by opening ends of Conway's normal form. (b) When  $\text{trunk}(K) \geq 6$  we can choose  $p$  and  $q$  so that each arc of  $\pi_K - \{p, q\}$  contains at least two critical points. (c) By pulling out a part, we can construct a local knot pattern.

the Conway's normal form is straight, we have a local knot pattern by opening both ends as in Figure 2(a).

By [1, Lemma 3(1)], from a knotted polygon in a  $2 \times 1$  tube with the smallest span we can obtain a proper knot pattern with the smallest span in the  $2 \times 1$  tube for that knot type by opening both ends of the polygon. For  $K \in \{3_1, 4_1, 5_1, 5_2\}$ , by applying the argument of [1, Theorem 4], we can completely characterise the configurations of  $K$  with smallest span, see Figures 3(a), (b), (c), (d) for examples. We can then conclude that the resulting proper knot patterns are all non-local. On the other hand, in these cases a local proper knot pattern can be constructed by increasing the span by one by the same method as in Figure 1(c).  $\square$

## References

- [1] K Ishihara, M Pouokam, A Suzuki, R Scharein, M Vazquez, J Arsuaga, and K Shimokawa, *Bounds for minimum step number of knots confined to tubes in the simple cubic lattice*, J. Phys A: Math. Theor. **50** (2017), 215601.
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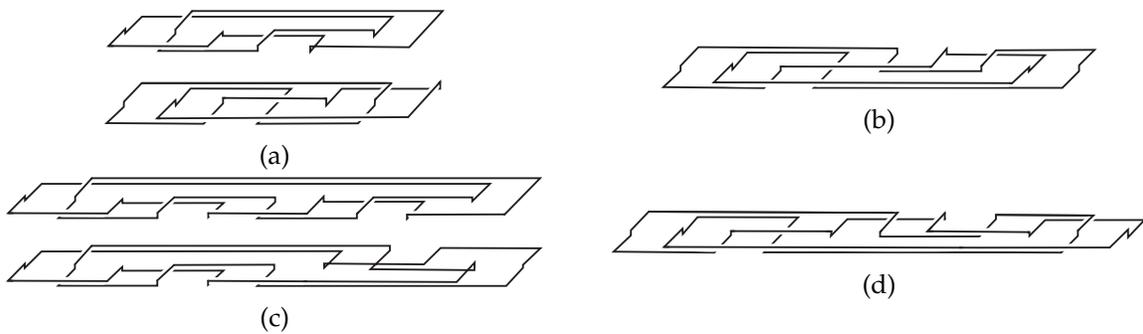


Figure 3: (a) Two polygons of  $3_1$  in  $2 \times 1$  tube with the smallest span 6; the first consists of 36 edges and the second consists of 38 edges. (b) A polygon of  $4_1$  in  $2 \times 1$  tube with the smallest span 8; this consists of 50 edges. (c) Two polygons of  $5_1$  in  $2 \times 1$  tube with the smallest span 10; these consist of 60 edges. (d) A polygon of  $5_2$  in  $2 \times 1$  tube with the smallest span 10; this consists of 62 edges.